

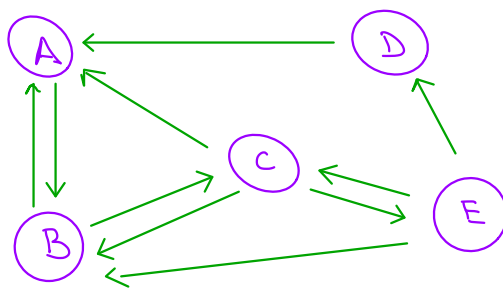
# 4.9 Markov Chains + Google's Page Rank Algorithm

Thanks to [kleinproject.org](http://kleinproject.org)

Goals:

- ① Rough intro to Google's Page Rank Alg.
- ② Intro to Markov chains

Let's explore the internet. There are pages and there are links from one page to another. Here is a simple example.



Q: How do we determine how "important" each page is so we can rank them.

Google's A: Take a random "walk", i.e. randomly click on links. Then...

more important page  $\iff$  you are more likely to end up on the page after several clicks

## The formal setup

- Record the probability of transitioning from one page to another in a transition matrix.

$$T = \begin{matrix} & \begin{matrix} A & B & C & D & E \end{matrix} \\ \begin{matrix} A \\ B \\ C \\ D \\ E \end{matrix} & \begin{bmatrix} 0 & \frac{1}{2} & \frac{1}{3} & 1 & 0 \\ 1 & 0 & \frac{1}{3} & 0 & \frac{1}{3} \\ 0 & \frac{1}{2} & 0 & 0 & \frac{1}{3} \\ 0 & 0 & 0 & 0 & \frac{1}{3} \\ 0 & 0 & \frac{1}{3} & 0 & 0 \end{bmatrix} \end{matrix}$$

prob. of moving from C to B

\*  $j$ th column contains the probabilities of moving from page  $j$  to page  $i$ .

- Record the probabilities that we are on each page after 0 clicks, 1 click, ... in a sequence of state vectors  $\bar{x}_0, \bar{x}_1, \bar{x}_2, \dots$ .

- Suppose we start on page C. Then

$$\bar{x}_0 = \begin{bmatrix} 0 \\ 0 \\ 1 \\ 0 \\ 0 \end{bmatrix} \begin{matrix} A \\ B \\ C \\ D \\ E \end{matrix}$$

- Looking at the graph and where we can go from page C, we see

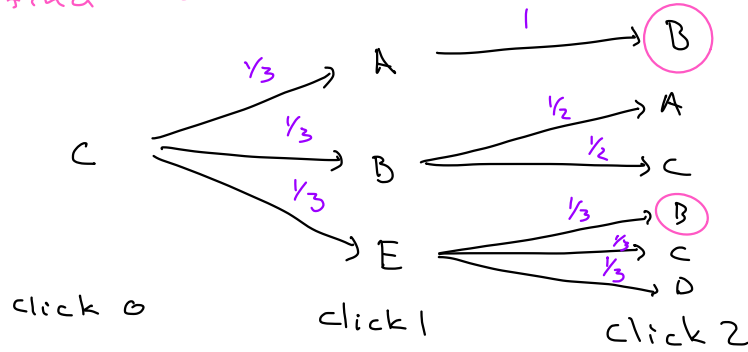
$$\bar{x}_1 = \begin{bmatrix} \frac{1}{3} \\ \frac{1}{3} \\ 0 \\ 0 \\ \frac{1}{3} \end{bmatrix}$$

• what is  $\bar{X}_2$ ?

$$\bar{X}_2 = \begin{bmatrix} * \\ * \\ * \\ * \\ * \end{bmatrix}$$

← prob of being at A after 2 clicks  
 ← " " " " B " 2 "

Let's find this one



$$\frac{1}{3} \cdot 1 + \frac{1}{3} \cdot \frac{1}{3} = \frac{4}{9}$$

... continue ... until you notice the following

③  $\bar{X}_k$  can be computed using T

$$\bar{X}_1 = T \cdot \bar{X}_0$$

$$\bar{X}_2 = T \cdot \bar{X}_1 = T^2 \bar{X}_0$$

$$\bar{X}_3 = T \cdot \bar{X}_2 = T^3 \bar{X}_0$$

⋮

$$\bar{X}_k = T \bar{X}_{k-1} \text{ so } \bar{X}_k = T^k \bar{X}_0$$

## Pause for definitions

Def A vector with nonnegative entries that add up to 1 is called a probability vector. A square matrix whose columns are probability vectors, is called a stochastic matrix

\* each state vector  $\bar{x}_0, \bar{x}_1, \dots$  is a probability vector.

\*  $T$  is a stochastic matrix.

Def A Markov chain is a sequence of prob. vectors together with a stochastic matrix  $T$  s.t.  $\bar{x}_k = T \bar{x}_{k-1}$  for all  $k \geq 1$ .

\* just like in our example.

## "Solving" the Page Rank Problem

we want to see what happens to our Markov chain after lots of clicks, i.e. as  $k \rightarrow \infty$ .

• Remember,  $\bar{x}_k = T^k \bar{x}_0$ , so we need to compute  $T^k$  for large  $k$  — yikes!

① Diagonalize  $T$  (if possible)

(a) Eigenvalues (using wolframalpha)

$$P(\lambda) = -\lambda^5 + \frac{7}{9}\lambda^3 + \frac{2}{9}\lambda^2 - \frac{1}{18}\lambda + \frac{1}{18}$$

There are 5 different eigenvalues, so T is diagonalizable. Let's call them

$\lambda_1, \lambda_2, \lambda_3, \lambda_4, \lambda_5$ . Then

$$\lambda_1 = 1$$

$$|\lambda_2| = |\lambda_3| \approx 0.7$$

$$|\lambda_4| = |\lambda_5| \approx 0.3$$

complex eigenvalues

(b) Let  $\bar{v}_i$  be an eigenvector assoc. to  $\lambda_i$ .

(we can compute these if needed)

$$\text{(e.g. } \bar{v}_1 = \begin{bmatrix} 4 \\ 5.3 \\ 3 \\ 0.3 \\ 1 \end{bmatrix} \text{)}$$

Then  $T = P D P^{-1}$ , where

$$D = \begin{bmatrix} 1 & & & & \\ & \lambda_2 & & & \\ & & \lambda_3 & & \\ & & & 0 & \\ & & & & \lambda_4 \\ & & 0 & & & \lambda_5 \end{bmatrix}, \quad P = [\bar{v}_1 \ \bar{v}_2 \ \bar{v}_3 \ \bar{v}_4 \ \bar{v}_5]$$

so

$$T^k = (P D P^{-1})^k = P D^k P^{-1}.$$

Now, we want to study  $T^k$  as  $k \rightarrow \infty$ .

$$T^\infty = \lim_{k \rightarrow \infty} T^k = \lim_{k \rightarrow \infty} P D^k P^{-1}$$

$$= \lim_{k \rightarrow \infty} P \begin{bmatrix} 1^k & & & & \\ & \lambda_2^k & & & \\ & & \lambda_3^k & & \\ & & & \lambda_4^k & \\ & & & & \lambda_5^k \end{bmatrix} P^{-1}$$

\* Recall:  $|\lambda_i| < 1$  for  $i=2,3,4,5$

$$= P \begin{bmatrix} 1 & & & & \\ & 0 & & & \\ & & \ddots & & \\ & & & 0 & \\ & & & & 0 \end{bmatrix} P^{-1}$$

\* compute  $P, P^{-1}$

$$\approx \begin{bmatrix} 0.293 & 0.293 & & & 0.293 \\ 0.390 & 0.390 & \dots & & 0.390 \\ 0.220 & 0.220 & & & 0.220 \\ 0.024 & 0.024 & & & 0.024 \\ 0.073 & 0.073 & & & 0.073 \end{bmatrix}$$

Thus, after lots of clicks we arrive at

$$\bar{x}_\infty = T^\infty \cdot \bar{x}_0 = T^\infty \begin{bmatrix} 0 \\ 0 \\ 1 \\ 0 \\ 0 \end{bmatrix} = \text{3rd col of } T^\infty = \begin{bmatrix} 0.293 \\ 0.390 \\ 0.220 \\ 0.024 \\ 0.073 \end{bmatrix}$$

Recall,  $\bar{x}_0 = \begin{bmatrix} 0 \\ 0 \\ 1 \\ 0 \\ 0 \end{bmatrix}$  meant we started on page C, but we see that we would have gotten the same answer no matter what page we started on.

The vector  $\bar{x}_\infty$  gives the prob. of ending up on each page after lots of clicks, no matter where we started. This gives the page ranking B, A, C, E, D

### Another view of what happened

Recall: the eigenvectors  $\bar{v}_1, \dots, \bar{v}_5$  are L.I.

so they form a basis for  $\mathbb{R}^5$  since they must also span. Then

$$\bar{x}_0 = c_1 \bar{v}_1 + \dots + c_5 \bar{v}_5 \quad \text{for some}$$

scalars  $c_1, \dots, c_5$ . Then,

$$\begin{aligned} T^k \bar{x}_0 &= c_1 T^k \bar{v}_1 + \dots + c_5 T^k \bar{v}_5 \\ &= c_1 \lambda_1^k \bar{v}_1 + \dots + c_5 \lambda_5^k \bar{v}_5 \end{aligned}$$

$$\begin{aligned} T \bar{v}_1 &= \lambda_1 \bar{v}_1 \\ T \bar{v}_2 &= \lambda_2 \bar{v}_2 \\ &\vdots \end{aligned}$$

so as  $k \rightarrow \infty$

$$\begin{aligned} T^k \bar{x}_0 &= c_1 1 \cdot \bar{v}_1 + c_2 0 \bar{v}_2 + \dots + c_5 0 \bar{v}_5 \\ &= c_1 \bar{v}_1 \end{aligned}$$

Also,  $c_1 \bar{v}_1$  will have to be a probability vector

so  $T^k \bar{x}_0$  is

- an eigen vector assoc. to  $\lambda=1$
- a vector whose entries sum to 1.

# Summary

Def If  $T$  is a stochastic matrix, then  $\bar{z}$  is a steady state vector for  $T$  if

- $T\bar{z} = \bar{z}$ , and
- $\bar{z}$  is a prob. vector.

Theorem If  $T$  is a stochastic matrix s.t. some power of  $T$  contains only positive entries, then  $T$  has a unique steady-state vector  $\bar{z}$ , and any Markov chain defined by  $\bar{x}_k = T\bar{x}_{k-1}$  converges to  $\bar{z}$  as  $k \rightarrow \infty$ , i.e.  $\bar{x}_\infty = \bar{z}$ .

"regular" ←

## Solving our Page rank problem (again)

① Create the transition matrix  $T$   
If it isn't regular, then tweak it.

our example: all entries in  $T^9$  are positive ✓

② Find any eigenvector associated to  $\lambda=1$ .  
(using a computer)

our example:

using wolframalpha

"eigenvector  $\{0, 1/2, 1/3, 1, 0\}, \{1, 0, \dots\}$ "

$$\bar{v}_1 = \begin{bmatrix} 4 \\ 5\frac{1}{3} \\ 3 \\ \frac{1}{3} \\ 1 \end{bmatrix}$$

③ Divide out the sum of entries to get the steady-state vector  $\bar{z}$ . Entries of  $\bar{z}$  give the page rank.

our example

sum entries in  $v_1$  is  $4\frac{1}{3}$

$$\bar{z} = \frac{3}{71} \bar{v}_1 \approx \begin{bmatrix} 0.293 \\ 0.390 \\ 0.220 \\ 0.024 \\ 0.073 \end{bmatrix}$$

← this gives the ranking

#1 #2 #3 #4 #5  
ⓑ ⓐ ⓒ ⓔ ⓓ